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Parallel Programming in C with MPI and OpenMP

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Chapter 13

Finite Difference Methods

Outline

Ordinary and partial differential equations
Finite difference methods
Vibrating string problem
Steady state heat distribution problem

Ordinary and Partial Differential Equations

- Ordinary differential equation: equation containing derivatives of a function of one variable
- Partial differential equation: equation containing derivatives of a function of two or more variables

Examples of Phenomena Modeled by PDEs

- Air flow over an aircraft wing
- Blood circulation in human body
- Water circulation in an ocean
- Bridge deformations as its carries traffic
- Evolution of a thunderstorm
- Oscillations of a skyscraper hit by earthquake
- Strength of a toy

Model of Sea Surface Temperature in Atlantic Ocean



Courtesy MICOM group at the Rosenstiel School of Marine and Atmospheric Science, University of Miami

Solving PDEs

Finite element method Finite difference method (our focus) Converts PDE into matrix equation • Result is usually a sparse matrix Matrix-based algorithms represent matrices explicitly Matrix-free algorithms represent matrix values implicitly (our focus)

Linear Second-order PDEs

Linear second-order PDEs are of the form Au_{xx} + 2Bu_{xy} + Cu_{yy} + Eu_x + Fu_y + Gu = H where A - H are functions of x and y only
Elliptic PDEs: B² - AC < 0
Parabolic PDEs: B² - AC = 0
Hyperbolic PDEs: B² - AC > 0 Copyright © The McGraw-Hill Companies, Inc. Permission required for reproduction or display.

Difference Quotients



Formulas for 1st, 2nd Derivatives

$$f'(x) = \frac{f(x+h/2) - f(x-h/2)}{h} + O(h^2)$$

f(x+h/2) = f(x-h/2)

$$f''(x) = \frac{f(x+h) - 2f(x) + f(x-h)}{h^2} + O(h^2)$$

Vibrating String Problem



Vibrating string modeled by a hyperbolic PDE

Solution Stored in 2-D Matrix

- Each row represents state of string at some point in time
- Each column shows how position of string at a particular point changes with time

Discrete Space, Time Intervals Lead to 2-D Matrix



Heart of Sequential Program

u[j+1][i] = 2.0*(1.0-L)*u[j][i] + L*(u[j][i+1] + u[j][i-1]) - u[j-1][i];



Parallel Program Design

- Associate primitive task with each element of matrix
- Examine communication pattern
- Agglomerate tasks in same column
- Static number of identical tasks
- Regular communication pattern
- Strategy: agglomerate columns, assign one block of columns to each task

Result of Agglomeration and Mapping



Communication Still Needed

- Initial values (in lowest row) are computed without communication
- Values in black cells cannot be computed without access to values held by other tasks









Ghost Points

Ghost points: memory locations used to store redundant copies of data held by neighboring processes

Allocating ghost points as extra columns simplifies parallel algorithm by allowing same loop to update all cells

Matrices Augmented with Ghost Points



Lilac cells are the ghost points.

Communication in an Iteration



This iteration the process is responsible for computing the values of the yellow cells.

Computation in an Iteration



This iteration the process is responsible for computing the values of the yellow cells. The striped cells are the ones accessed as the yellow cell values are computed.

Complexity Analysis

- Computation time per element is constant, so sequential time complexity per iteration is $\Theta(n)$
- Elements divided evenly among processes, so parallel computational complexity per iteration is $\Theta(n / p)$
- During each iteration a process with an interior block sends two messages and receives two messages, so communication complexity per iteration is Θ(1)

Isoefficiency Analysis

- Sequential time complexity: $\Theta(n)$
- Parallel overhead: $\Theta(p)$
- Isoefficiency relation:
 - $n \ge Cp$
- To maintain the same level of efficiency, n must increase at the same rate as p
- If $M(n) = n^2$, algorithm has poor scalability
- If matrix of 3 rows rather than *m* rows is used, M(n) = n and system is perfectly scalable

Replicating Computations

- If only one value transmitted, communication time dominated by message latency
- We can reduce number of communications by replicating computations
- If we send two values instead of one, we can advance simulation two time steps before another communication

Replicating Computations

Without replication:



With replication:



Communication Time vs. Number of Ghost Points



Ghost Points

Steady State Heat Distribution Problem

Ice bath



Solving the Problem

Underlying PDE is the Poisson equation

$$u_{xx} + u_{yy} = f(x, y)$$

This is an example of an elliptical PDE
Will create a 2-D grid
Each grid point represents value of state solution at particular (*x*, *y*) location in plate

Heart of Sequential Program

w[i][j] = (u[i-1][j] + u[i+1][j] + u[i][j-1] + u[i][j+1]) / 4.0;



Parallel Algorithm 1

- Associate primitive task with each matrix element
- Agglomerate tasks in contiguous rows (rowwise block striped decomposition)
- Add rows of ghost points above and below rectangular region controlled by process

Ghost Points

Ghost points: memory locations used to store redundant copies of data held by neighboring processes

Allocating ghost points as extra rows simplifies parallel algorithm by allowing same loop to update all cells

Example Decomposition









16 × 16 grid divided among 4 processors

Complexity Analysis

Sequential time complexity: $\Theta(n^2)$ each iteration Parallel computational complexity: $\Theta(n^2 / p)$ each iteration Parallel communication complexity: $\Theta(n)$ each iteration (two sends and two receives of *n* elements)

Isoefficiency Analysis

 Sequential time complexity: Θ(n²)
 Parallel overhead: Θ(pn)
 Isoefficiency relation: n² ≥ Cnp ⇒ n ≥ Cp

$$M(Cp) / p = C^2 p^2 / p = C^2 p$$

This implementation has poor scalability

Parallel Algorithm 2

- Associate primitive task with each matrix element
- Agglomerate tasks into blocks that are as square as possible (checkerboard block decomposition)
- Add ghost points to all four sides of rectangular region controlled by process

Example Decomposition



16 × 16 griddivided among16 processors

Implementation Details

- Using ghost points around 2-D blocks requires extra copying steps
- Ghost points for left and right sides are not in contiguous memory locations
- An auxiliary buffer must be used when receiving these ghost point values
- Similarly, buffer must be used when sending column of values to a neighboring process

Complexity Analysis

Sequential time complexity: $\Theta(n^2)$ each iteration Parallel computational complexity: $\Theta(n^2 / p)$ each iteration Parallel communication complexity: $\Theta(n/\sqrt{p})$ each iteration (four sends and four receives of n / \sqrt{p} elements each)

Isoefficiency Analysis

Sequential time complexity: Θ(n²)
Parallel overhead: Θ(n √p)
Isoefficiency relation: n² ≥ Cn √p ⇒ n ≥ C √p

 $M(C_{\sqrt{p}})/p = C^2 p/p = C^2$

This system is perfectly scalable

Gauss-Seidel: Red-black ordering



"Red-black" GS iteration

- 1. Exchange black edge values with neighboring processes
- 2. Perform GS updates on all red points
- 3. Exchange red edge values with neighboring processes
- 4. Perform GS updates on all black points

Summary (1/4)

- PDEs used to model behavior of a wide variety of physical systems
- Realistic problems yield PDEs too difficult to solve analytically, so scientists solve them numerically
- Two most common numerical techniques for solving PDEs
 - finite element method
 - finite difference method

Summary (2/4)

Finite different methods
 Matrix-based methods store matrix explicitly

- Matrix-free implementations store matrix implicitly
- We have designed and analyzed parallel algorithms based on matrix-free implementations

Summary (3/4)

Linear second-order PDEs
Elliptic (e.g., heat equation)
Hyperbolic
Parabolic (e.g., wave equation)
Hyperbolic PDEs typically solved by methods not as amenable to parallelization

Summary (4/4)

- Ghost points store copies of values held by other processes
- Explored increasing number of ghost points and replicating computation in order to reduce number of message exchanges
- Optimal number of ghost points depends on characteristics of parallel system